



# JSR 292 Cookbook

Rémi Forax

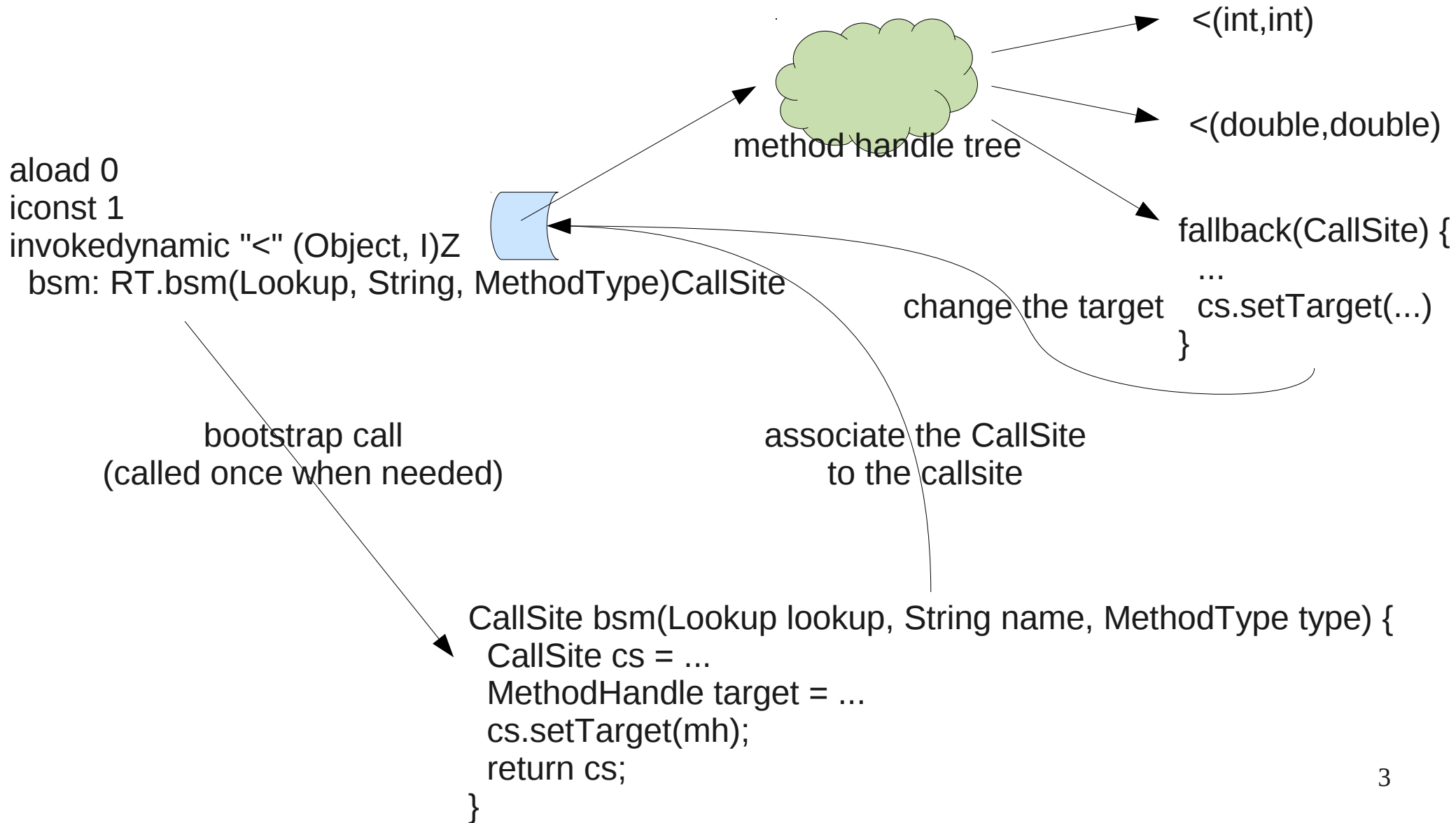
# Goals

How to use JSR 292 to create common dynamic language runtime patterns

Try to gather best practices

Exercise the JSR 292 spec/implementations

# invokedynamic



# JSR 292

enhanced bytecode + a new API

new constant pool constants

invokedynamic

link a callsite to one or several target method implementations ?

can relink dynamically !

java.lang.invoke

manage function pointers (MethodHandle)

combinators

provide adhoc classes: ClassValue, SwitchPoint

# JSR 292 & Java

JSR 292 is **poorly supported** by Java  
(the language)

No support of

invokedynamic

expando keyword ? => expando type, expando method

constant method handle

Constant method reference (lambda) should be  
converted to a MethodHandle (not currently specified)

# Cookbook

interceptors

constants lazy initialization

callsite adaptation

varargs, spread, named parameters

method dispatch

single-dispatch

monomorphic IC, unverified entry point, bi-morphic IC, dispatch table

double-dispatch

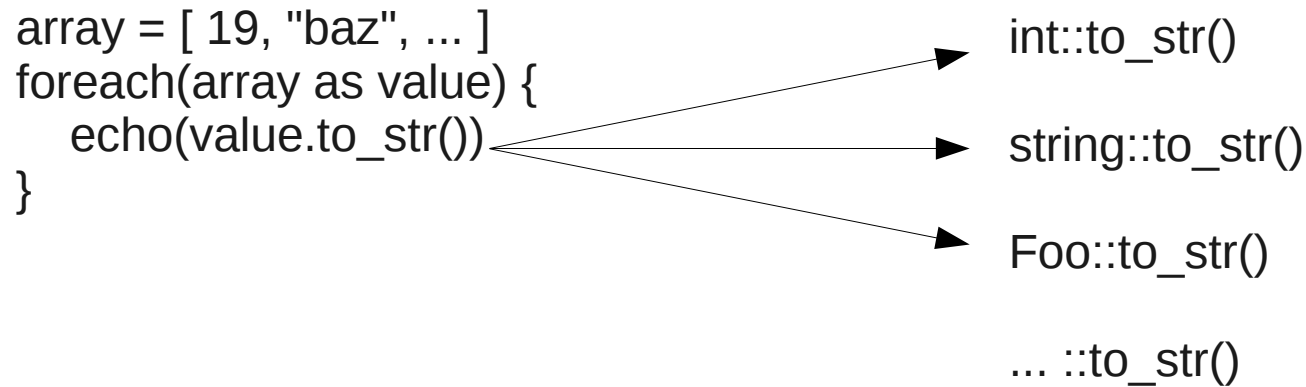
binary op, visitors

multi-dispatch

metaclass & invalidation

# Single Dispatch

The target method is chosen depending on the class of the receiver



Can't use a vtable !

Constructs a dispatch table ?

One by callsite or one by selector

# Dispatch Table

Hash map between a class and a method handle

Use invoker + fold to insert the target in front of the arguments

```
DispatchMap dispatchMap = new DispatchMap() {
    protected MethodHandle findMethodHandle(Class<?> receiverClass) {
        MethodHandle target = ...
        return target.asType(type);
    }
};
MethodHandle lookupMH = MethodHandles.filterReturnValue(Object#getClass,
    DispatchMap#lookup.bindTo(dispatchMap));
lookupMH = lookupMH.asType(methodType(MethodHandle.class, type.parameterType(0)));
MethodHandle target = MethodHandles.foldArguments(
    MethodHandles.exactInvoker(type), lookupMH);
callsite.setTarget(target);
```

```
class DispatchMap {
    public MethodHandle lookup(Class<?> k) {
```



# Perf ? How it works ?

If a method handle is

- Used in a hot code

- Static or reachable from a static context

  - invokedynamic callsite

  - static final method handle

  - constant method handle (+ldc)

The JIT will inline the whole method handle blob at callsite

- Gold card: don't decrease inlining\_depth

- Others thresholds still exist (number of IR nodes, etc)

# Perf ? Dispatch Table ?

So the method handles get from the dispatch table **aren't inlined** !

Use an inlining cache !

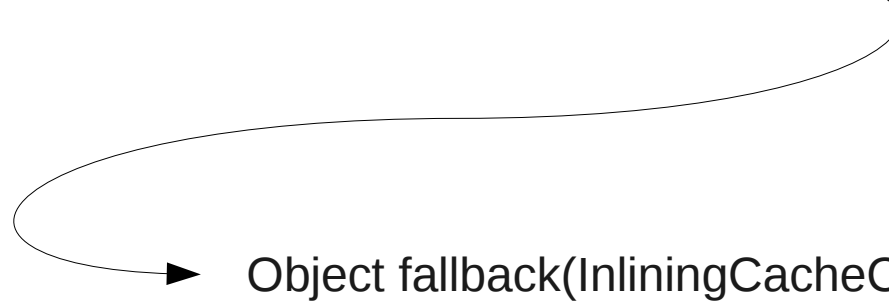
Keep last returned MH and check if the receiver class has not changed

Constructs a tree of decision and fallback to a dispatch table if depth > threshold

# Inlining cache

Delay the computation of the target until arguments are available

```
array = [ 19, "baz", 42.0 ]  
foreach(array as value) {  
  echo(value.to_str())  
}
```



# Inlining cache

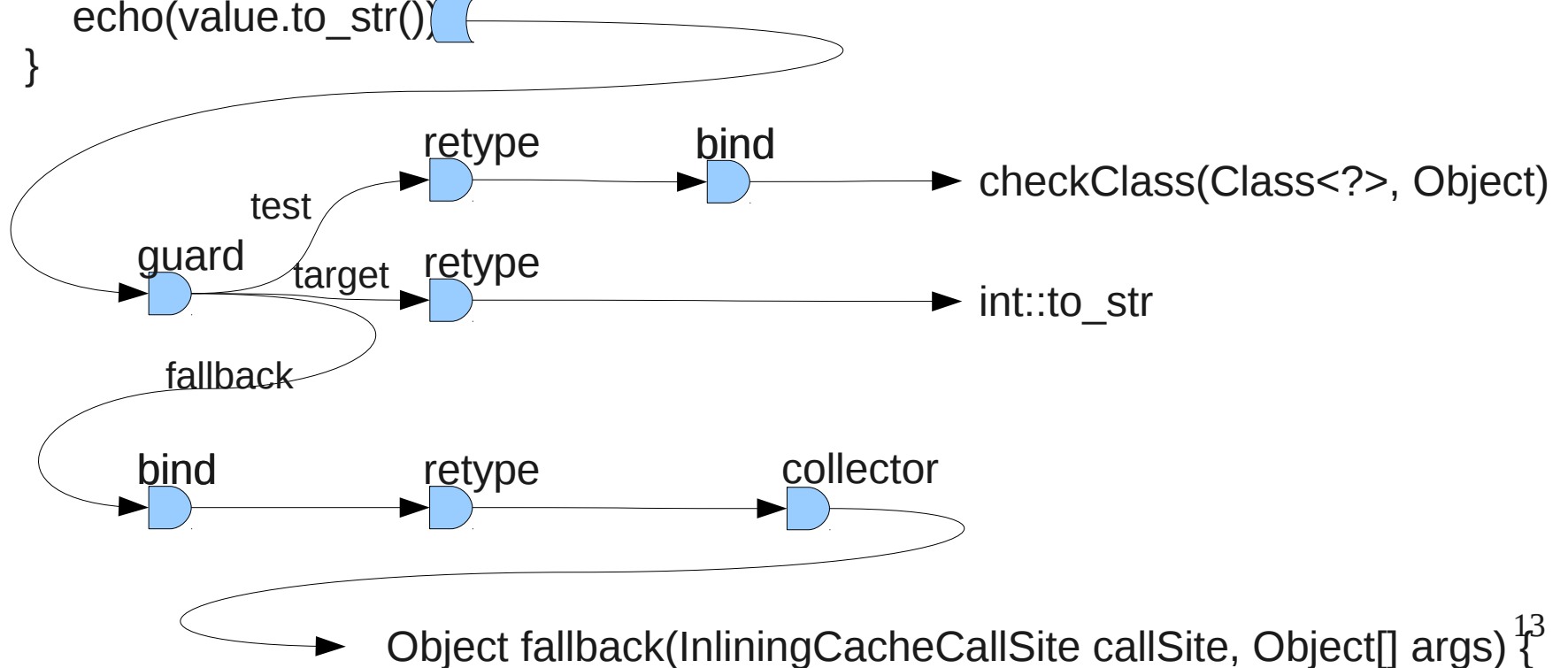
The BSM installs fallback as target

```
class InliningCacheCallSite extends MutableCallSite {  
    ...  
}  
  
static CallSite bootstrap(Lookup lookup, String name, MethodType type) {  
    InliningCacheCallSite callSite = new InliningCacheCallSite(...);  
  
    MethodHandle fallback = #fallback.bindTo(callSite);  
    fallback = fallback.asCollector(Object[].class, type.parameterCount());  
    fallback = fallback.asType(type);  
  
    callSite.setTarget(fallback);  
    return callSite;  
}  
  
static Object fallback(InliningCacheCallSite callSite, Object[] args) {
```

# Inlining cache

Install a guard to avoid to do the lookup each time

```
array = [ 19, "baz", 42.0 ]  
foreach(array as value) {  
  echo(value.to_str())  
}
```



# Inlining cache

The guard fallback reuse the previous target

```
Object fallback(InliningCacheCallSite callSite, Object[] args) throws Throwable {
    MethodType type = callSite.type();
    Class<?> receiverClass = args[0].getClass();
    MethodHandle target = ...
    target = target.asType(type);
    MethodHandle test = #checkClass.bindTo(receiverClass);
    test = test.asType(test.type().changeParameterType(0, type.parameterType(0)));
    MethodHandle guard = MethodHandles.guardWithTest(test, target, callSite.getTarget());
    callSite.setTarget(guard);
    return target.invokeWithArguments(args);
}

static boolean checkClass(Class<?> clazz, Object receiver) {
    return receiver.getClass() == clazz;
}
```

# Inlining cache

A chain of guards too big will kill performance

Store the depth in the CallSite object

```
Object fallback(InliningCacheCallSite callSite, Object[] args) throws Throwable {
    MethodType type = callSite.type();
    if (callSite.depth >= MAX_DEPTH) {
        ...
    }
    Class<?> receiverClass = args[0].getClass();
    MethodHandle target = ...
    target = target.asType(type);
    MethodHandle test = #checkClass.bindTo(receiverClass);
    test = test.asType(test.type().changeParameterType(0, type.parameterType(0)));
    MethodHandle guard = MethodHandles.guardWithTest(test, target, callSite.getTarget());
    callSite.depth++;
    callSite.setTarget(guard);
    return target.invokeWithArguments(args);
}
```

# Inlining cache – Thread safety

The code is not thread safe

race to call setTarget

no problem it's a cache

depth not atomic !

real depth may be greater than MAX\_DEPTH

in slow path, so better to use volatile + CAS (AtomicInteger)

```
Object fallback(InliningCacheCallSite callSite, Object[] args) throws Throwable {
    MethodType type = callSite.type();
    if (callSite.depth.get() >= MAX_DEPTH) {
        ...
    }
    ...
    callSite.depth.incrementAndGet();
    callSite.setTarget(guard);
    return target.invokeWithArguments(args);
}
```



# How to improve the dispatch table ?

Solution if few method handles:

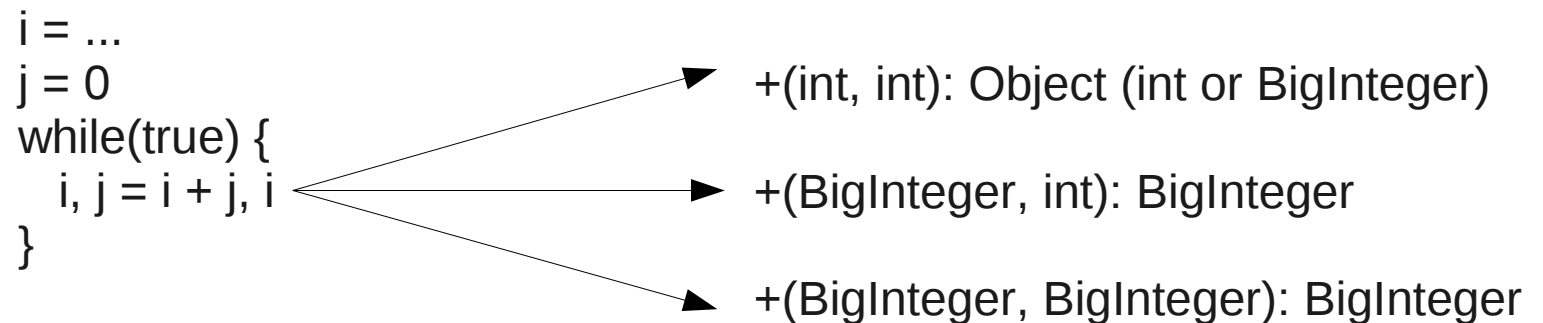
- Dispatch map (class -> int) + **Switch inliner**
  - switchInliner(MH fallback, MH... mhs)
    - All method handles must have the same signature
    - Returns a mh that takes a supplementary first parameter (the index) and calls mhs[index].

The JIT should try to inline the all mhs

- Also solve problem of several class that share the same method implementation (inheritance)

# Binary operations

Double dispatch, the target method depends on the class of the two arguments



Moreover, int operations can overflow to BigInteger

```

Object fallbackOpBoth(BinOpCallsite callSite, Object value1, Object value2) {
    Class<?> class1 = value1.getClass(), class2 = value2.getClass();
    MethodHandle target, guard1, guard2;
    if (class1 == BigInteger.class) {
        guard1 = BIGINTEGER_CHECK;
        if (class2 == BigInteger.class) {
            guard2 = BIGINTEGER_CHECK2;
            target = BigInteger#add;
        } else {
            if (class2 != Integer.class) { throw ... }
            guard2 = INTEGER_CHECK2;
            target = MethodHandles.filterArguments(
                BigInteger#add, 1, OBJECT_TO_INTEGER_TO_BIGINTEGER);
        }
    } else {
        ...
    }
    target = target.asType(callSite.type());
    MethodHandle fallback = callSite.getTarget();
    MethodHandle guard = MethodHandles.guardWithTest(guard1,
        MethodHandles.guardWithTest(guard2, target, fallback),
        fallback);
    callSite.setTarget(guard);
    return target.invoke(value1, value2);
}

```

# Perf warning: unboxing trouble !

Conversion Object -> int is not equivalent to  
Object -> Integer -> int

```
mh = BigInteger#valueOf(long);  
mh = mh.asType(methodType(BigInteger.class, int.class));  
mh = mh.asType(methodType(BigInteger.class, Integer.class));  
mh = mh.asType(methodType(BigInteger.class, Object.class));  
OBJECT_TO_INTEGER_TO_BIGINTEGER = mh
```

Object -> int accepts Byte -> byte -> int

# Binary operations

Signature is fixed, 2 parameters, fixed small number of classes

=> no dispatch table needed

A lot of operation involve one constant

$x + 1$ ,  $1 + x$ , etc

=> no need to do a double dispatch

Usually overflow are rare

=> no need to construct the whole tree of possibility

If one arg is constant

=> can use dedicated overflow test !

# Overflow test for +

Check if values as same sign and result an opposite sign

```
Object safeAdd(int value1, int value2) {  
    int result = value1 + value2;  
    if ((value1 ^ result) < 0 && (value1 ^ value2) >= 0) {  
        return BigInteger.valueOf(value1).add(BigInteger.valueOf(value2));  
    }  
    return result;  
}
```

It can be simplified, the overflow test for  $x + 1$  is  
if ( $x == \text{Integer.MAX\_VALUE}$ )

# Overflow for $x + cst$

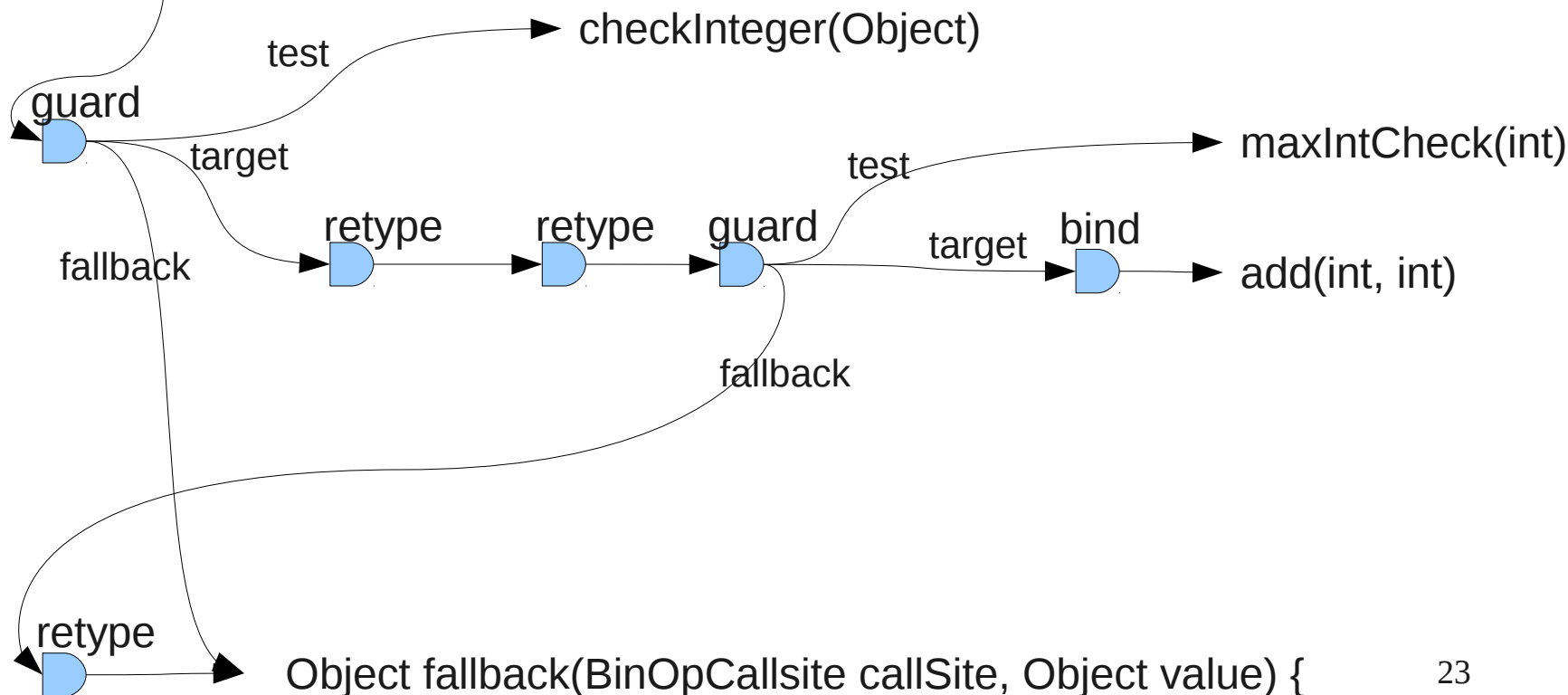
Reuse the same fallback for the guards

```
i = 0
```

```
while(true) {
```

```
  i = i + 1
```

```
}
```



# Overflow for $x + cst$

invokedynamic with one argument,  
constant value is sent as a bootstrap argument

```
Object fallbackOpLeft(BinOpCallsite callSite, Object value) {
    ...
    MethodHandle op = MethodHandles.insertArgument(#add(int, int), 1, callSite.value);
    MethodHandle overflowTest = createAddOverflowTest(rightValue);
    if (overflowTest != null) {
        overflowGuard = MethodHandles.guardWithTest(overflowTest,
            op,
            fallback.asType(methodType(Object.class, int.class)));
    } else {
        overflowGuard = op;
    }
    overflowGuard = overflowGuard.asType(methodType(Object.class, Integer.class)).
        asType(methodType(Object.class, Object.class));
    MethodHandle target = MethodHandles.guardWithTest(INTEGER_CHECK,
        overflowGuard, fallback);
}
```



# VM Optimization ?

*x binop* cst optimization enables the VM to do further optimizations

```
i = 0
while(i < array.length) {
    echo(array[i])
    i = i + 1
}
```

The VM should fold the 3 Integer checks  
then range check optimization can be applied  
then escape analysis remove boxing  
=> same loop as Java :)

# Metaclass & invalidation

The metaclass provides language specific class metadata

Metaclass can be mutable

```
MethodHandle mh = lookup.findVirtual(String.class, "toLowerCase",
    methodType(String.class));
for(int i=0; i<10; i++) {
    System.out.println("Hello".toUpperCase()); // invokedynamic
    if (i == 3) {
        MetaClass.getMetaClass(String.class).
            redirect("toUpperCase", methodType(String.class), mh);
    }
}
```

# ClassValue

Association between a class object and a runtime specific object

Act has a concurrent weak hash map

```
public class MetaClass {  
    private final HashMap<Selector, MethodHandle> vtable =  
        new HashMap<Selector, MethodHandle>();  
  
    private static final ClassValue<MetaClass> metaClassValue =  
        new ClassValue<MetaClass>() {  
            protected MetaClass computeValue(Class<?> type) {  
                return new MetaClass(...);  
            }  
        };  
  
    public static MetaClass getMetaClass(Class<?> clazz) {  
        return metaClassValue.get(clazz);  
    }  
}
```

# Managing mutation

## Pull

Volatile serial number in MetaClass

Bind the serial current value in the method handle blob

Bind the metaclass + field accessor

Increment when metaclass is mutated

## Push

1 frozen MetaClass <--> 1 SwitchPoint

Bind the SwitchPoint

Volatile boolean when interpreted

No check when JITed + dependencies for invalidation

Create a new SwitchPoint if mutated

# SwitchPoint

Insert a SwitchPoint check after the class check and before the target adaptation

```
Object staticFallback(InvokeStaticCallSite cs) {
    MetaClass metaClass = MetaClass.getMetaClass(cs.ownerType);
    MethodType type = cs.type();
    MethodHandle mh;
    SwitchPoint switchPoint;
    synchronized(MetaClass.MUTATION_LOCK) {
        mh = metaClass.staticLookup(cs.name, type);
        switchPoint = metaClass.switchPoint;
    }
    if (mh == null) { mh = ... }
    mh = mh.asType(type);
    MethodHandle target = switchPoint.guardWithTest(mh, fallback);
    ...
}
```

# SwitchPoint in HotSpot

SwitchPoint is currently implemented with a volatile check but

**Pull then push** optimization:

Detects that SwitchPoint is hot (specific profiling)

Make it a no-op + store callsites using a switch point in the switchPoint + de-optimization if invalidation

=> as fast as an inlining cache call !

# Invalidation & inheritance

Need to maintain a metaclass hierarchy to also invalidate sub-metaclasses

```
private static final ClassValue<MetaClass> metaClassValue =
    new ClassValue<MetaClass>() {
        protected MetaClass computeValue(Class<?> type) {
            Class<?> superclass = type.getSuperclass();
            MetaClass parentMetaClass = (superclass == null)? null: getMetaClass(superclass);
            return new MetaClass(parentMetaClass);
        }
    };
private final LinkedList<WeakReference<MetaClass>> subMetaClasses =
    new LinkedList<>();
MetaClass(MetaClass parent) {
    synchronized(MUTATION_LOCK) {
        switchPoint = new SwitchPoint();
        this.parent = parent;
        if (parent != null) {
            parent.subMetaClasses.add(new WeakReference<MetaClass>(this));
        }
    }
}
```

# Bulk invalidation

To try to avoid deoptimization flood

```
public void redirect(String name, MethodType type, MethodHandle target) {
    synchronized(MUTATION_LOCK) {
        ArrayList<SwitchPoint> switchPoints = new ArrayList<>();
        mutateSwitchPoints(this, switchPoints);
        SwitchPoint.invalidateAll(switchPoints.toArray(new SwitchPoint[switchPoints.size()]));
        vtable.put(new Selector(name, type), target);
    }
}

private static void mutateSwitchPoints(MetaClass mc, List<SwitchPoint> switchPoints) {
    switchPoints.add(mc.switchPoint);
    mc.switchPoint = new SwitchPoint();
    for(Iterator<WeakReference<MetaClass>> it = mc.subMetaClasses.iterator(); it.hasNext();) {
        MetaClass subMC = it.next().get();
        if (subMC == null) { it.remove(); continue; }
        mutateSwitchPoints(subMC, switchPoints);
    }
}
```



# Metaclass pattern & prototype

This pattern doesn't work with prototype based (not class based) languages

Self, JavaScript, Sesh, etc.

## Solutions ?

Use pseudo class trick (V8)

JVM will require two guards instead of 1

All Objects (even String) implement get/setMetaClass()

interface injection?

Use profiling + allocation site class

# Overall recommendations

Compiler should have a type inference pass

and check dead-code/missing return

Segregate fast path/slow path

No second class citizen

No RubyObject, GroovyObject, etc.

Design with concurrency in mind

# Questions ?



<https://code.google.com/p/jsr292-cookbook/>