Fast Bytecodes for Funny Languages

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THE #1 PLATFOR







- JVMs are used for lots of non-Javac bytecodes
- Bytecode patterns are different
- Bytecode producer (this crowd) is curious:
 - Are these bytecodes "fast"? "fast enough"?
 - Am I inlining?
 - JIT'ing to good code? (JIT'ing at all?)
- Personal goal: you ARE using that fancy JIT, right?
 - Otherwise, why did I bother? ;-)
- 'nother goal: maybe help the world escape the Java box
- New language must have a fast reputation
 - Or a large programmer population won't look

Can Your Language Go "To The Metal"?



- Can your bytecodes go "To The Metal"?
 - e.g. Can simple code be mapped to simple machine ops
- Methodology:
 - Write dumb hot SIMPLE loop in lots of languages
 - Look at performance
 - Look at JIT'd code
 - Look for mismatch between language & JVM & machine code
- NOT-GOAL: Who is fastest
- NOT TESTED:
 - Ease of programming, time-to-market, maintenance cost
 - Domain fit (e.g. Ruby-on-Rails)
- IGNORING:
 - Blatant language / microbenchmark mismatch (FixNum/BigNum)

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See "http://blogs.azulsystems.com/cliff/2008/09/a-plea-for-prog.html"

```
for( int i=1; i<10000000; i++ )
sum += (sum^i)/i;</pre>
```

- Run with Java, Scala, Clojure, JRuby, JPC, Javascript/Rhino
 - Willing to try more languages
- NOT Goal:
 - Discover the fastest
- NOT Fair:
 - I changed benchmark to be more "fair" to various languages
 - e.g. Using 'double' instead of 'int' for Javascript.
- Tested on Azul JVM
 - Because it's got the best low-level JVM perf analysis tools I've seen





• Java (unrolled, showing 1 iter):

- add rTmp0, rI, 5 // for unrolled iter #5
- xor rTmp1, rSum, rTmp0
- div rTmp2, rTmp1, rTmp0
- add rSum, rTmp2, rSum
- ... repeat for unrolled iterations
- Scala Same as Java!
 - Scala "Elegant Scala-ized version" Ugh
- Clojure Almost close; very allocation heavy; oddball 'holes'
- JRuby Major inlining not happening,
 - misled by debugging flag?
- JPC Fascinating; totally inlined X86 emulation
 - But JIT doesn't grok e.g. dead X86 flag setting
- Javascript/Rhino Death by 'Double' (not 'double')



- Java Unfair advantage. Semantics match JIT expectations.
- Scala Borrowed heavily from Java semantics
 - Close fit allows close translation
 - And thus nearly direct translation to machine code
 - Penalty: Only have Java semantics, e.g. no 'unsigned' type, no auto-BigNum inflation
- JPC Java "PC"; pure Java X86/DOS emulator
 - Massive bytecode generation; LOTS of JIT'ing
 - For this example: 16000 classes, 7800 compilations
 - But JIT falls down removing e.g. redundant X86 flag sets
 - Maybe fix by never storing 'flags' into memory
 - Also no fixnum issue





- Clojure "almost close"
 - Good: no obvious subroutine calls in inner loop
 - Bad: Massive "ephemeral" object allocation requires good GC
 - But needs Escape Analysis to go fast
 - Ugly: fix-num overflow checks everywhere
 - Can turn off fix-nums; could be same speed as Java
 - Weird "holes" Not-optimized reflection calls here & there
- Jython "almost close"
 - Also has Fixnum issue; massive allocation
 - Some extra locks thrown in
- JavaScript/Rhino
 - All things are Doubles not 'double'
 - Same allocation issues as Fixnum
 - Otherwise looks pretty good (no fixnum checks)





- Common Issue FixNums
 - Allocation costs (but GC does not); final fields cost mfence
 - Could do much better w/JIT
 - Need ultra-stupid Escape Analysis
 - Need some (easy) JIT tweaking
 - e.g. Getting around Integer.valueOf(i) caching
- JRuby Missed The Metal
 - Assuming CachingCallSite::call inlines (and allows further inlining)
 - Using +PrintInlining to determine
 - But flag is lying: claims inlined, but it's not
 - (issue is w/BimorphicInlining 'guessing' wrong target)
 - Confirmed w/debug Java5 & GDB on product-mode Java6
 - Confirmed w/Azul My 1st impression: can't be inlined ever
 - (but please Charles tell me why you think it should!)



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- CAN take "funny language" to the metal (e.g. Scala)
- Easy to be misled (JRuby)
 - Reliance on non-QA'd debugging flag +PrintInlining
 - Rules on inlining are complex, subtle
- But so much performance depends on it!
 - And exact details matter: class hierarchy, final, interface, singletarget
- And how do you know about JRockit (BEA), IBM, etc?
- And heuristics change anyways,
 - what works today is dog-slow tomorrow
 - Unless your language hits the standard benchmark list
- Language/bytecode mismatch made worse by assuming
 - e.g. Uber GC or Uber Escape Analysis, or subtle final-field semantics

- How do you when your bytecodes are working well?
 - Lack of good tools support for telling **why** optimizer bailed out
 - Same lack for C++, Fortran for the past 50 yrs
- 1- Run microbenchmark w/hot simple code & GDB it
 - Why? So you can 'read' the asm
 - Break in w/GDB and read the asm
 - Did translation go as you expected? (no?)
 - Simplify even more...
- 2- Run Your Fav JVM w/debugging flags
 - But debugging flags can lie
 - Must cross-correlate with (1)
- 3- Run on Azul & use RTPM
 - Only sorta kidding: email me & ask for Academic Account

Round 2: Alternative Concurrency

- Only Clojure looked at got lazy busy
- Clojure Traveling Salesmen Problem w/worker 'ants'
- Tried up to 600 'ants' on a 768-way Azul

Good scaling (but 20Gig/sec allocation)

- Tried contention microbenchmark
 - Performance died
 - Less TOTAL throughput as more CPUs added
 - JDK 6 Library failure? Clojure usage model failure?
 - Not graceful fall-back under pressure

- Too little data (e.g. no Scala), just my speculation
- Clojure-style MIGHT work, might not
 - 600-way thread-pool works well on Java also
- NOT graceful under pressure
 - Adding more CPUs should not be bad
 - Thru-put cap expected, or maybe slow degrade but rapid falloff!?!?
 - Maybe STM becomes "graceful under pressure" later
 - (but it's been 15yrs at still not good)
- Note: complex locking schemes suffer same way
 - eg. add more concurrent DB requests, DB throughput goes up
 - ..then down, then craters
 - Why does DB not queue requests & maintain max throughput?

Future Big Problem

- Reliable performance "under pressure"
 - Eg adding more Threads to a hot lock drops throughput some
 - ...then stabilizes throughput as more threads added
 - Eg adding more CPUs to a wait/notifyall peaks throughput
 - ...then falls constant as most threads uselessly wakeup & sleep
- And we're working w/weak & immature concurrency libs
- Everybody has a max-throughput
 - And Fall-off Under Load is Bad (FULB[™])
- Naively: just queue requests beyond saturation point
 - And maintain max-throughput
- Then why not just publish that (now reliable) max throughput
- So can predict performance as the min of max-thruput's...